

# Mobile ad hoc networks

Standard Mobile IP needs an infrastructure

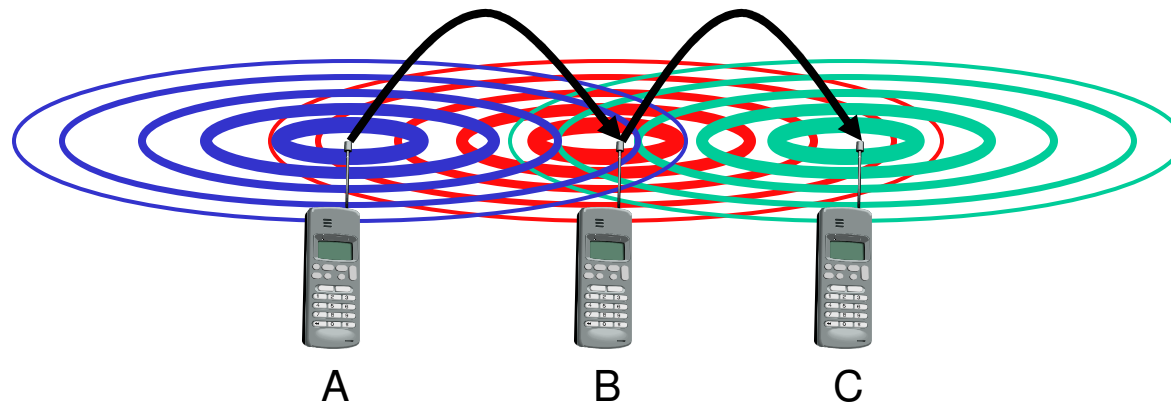
- ❑ Home Agent/Foreign Agent in the fixed network
- ❑ DNS, routing etc. are not designed for mobility

Sometimes there is no infrastructure!

- ❑ remote areas, ad-hoc meetings, disaster areas
- ❑ cost can also be an argument against an infrastructure!

Main topic: routing

- ❑ no default router available
- ❑ every node should be able to forward



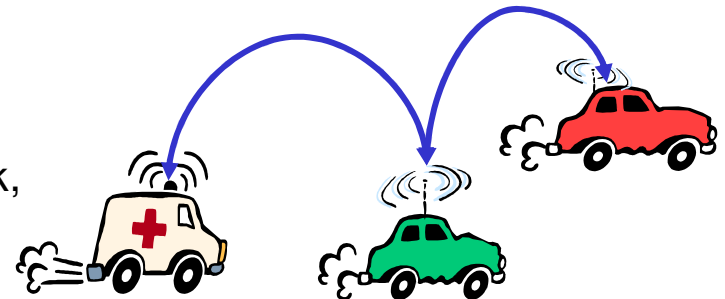
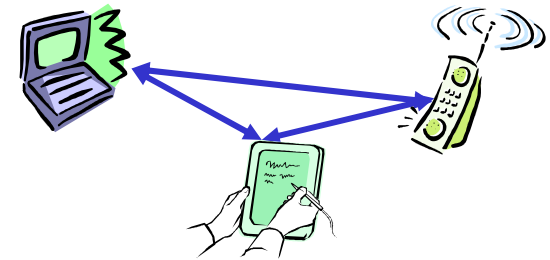
# Solution: Wireless ad-hoc networks

## Network without infrastructure

- ❑ Use components of participants for networking

## Examples

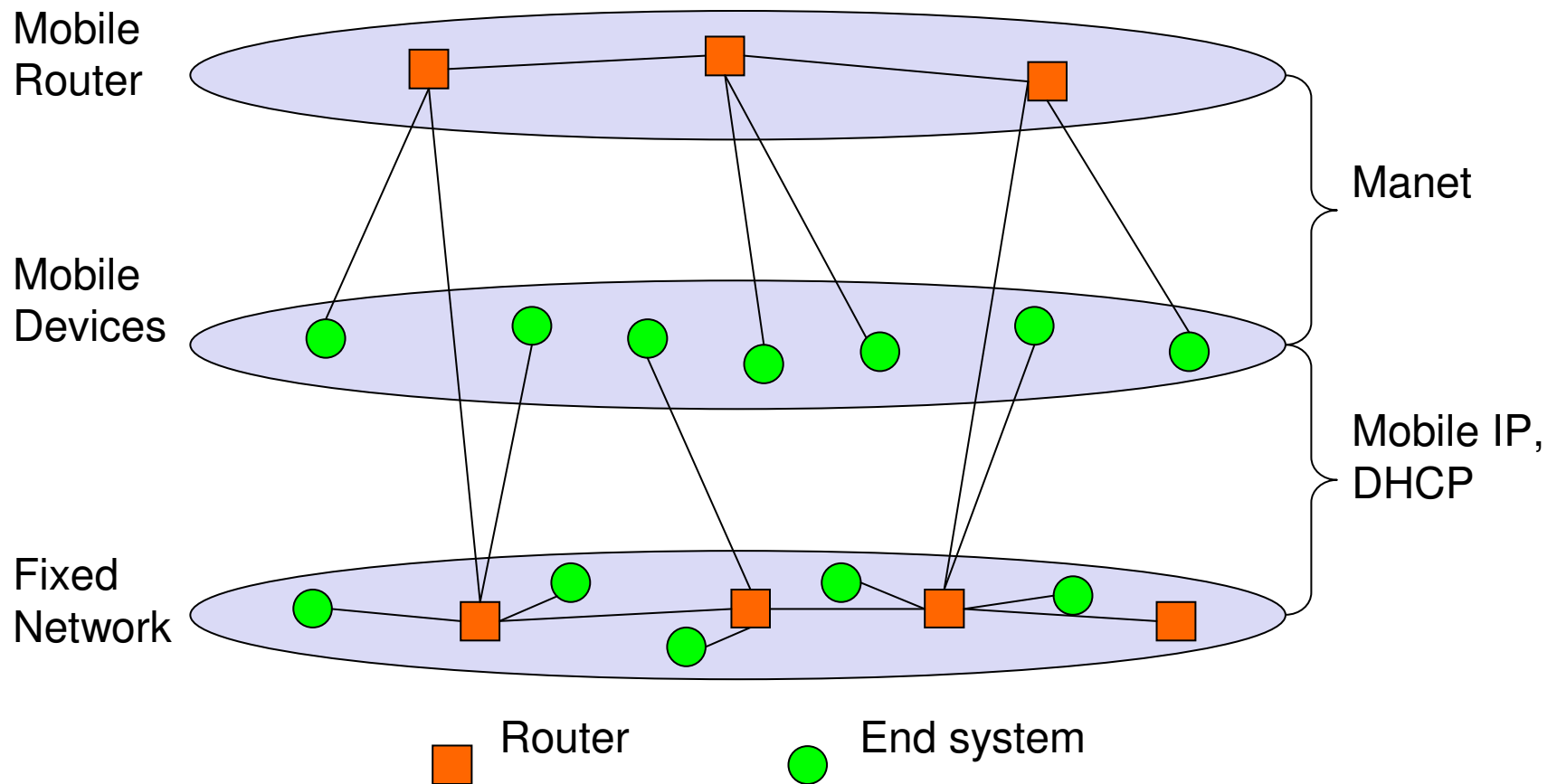
- ❑ Single-hop: All partners max. one hop apart
  - Bluetooth piconet, PDAs in a room, gaming devices...
- ❑ Multi-hop: Cover larger distances, circumvent obstacles
  - Bluetooth scatternet, TETRA police network, car-to-car networks...



## Internet: MANET (Mobile Ad-hoc Networking) group



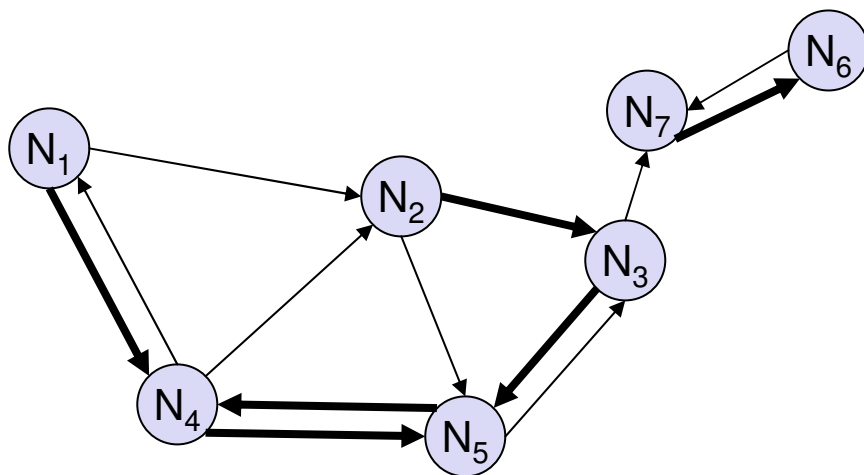
# Manet: Mobile Ad-hoc Networking



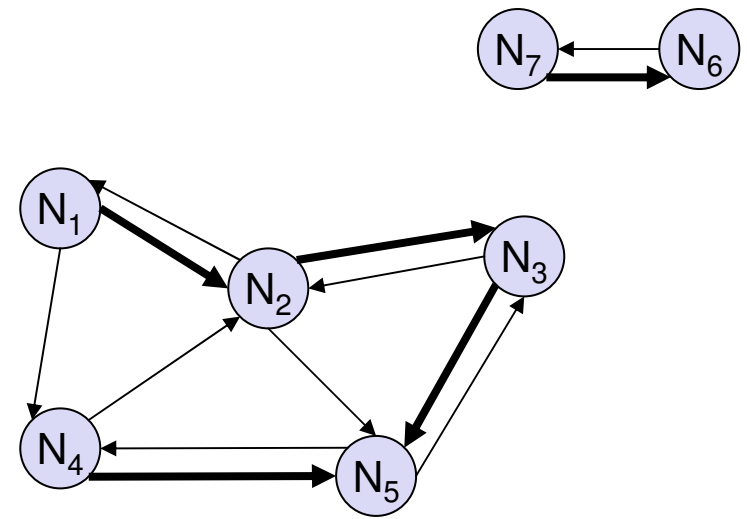
# Problem No. 1: Routing

## Highly dynamic network topology

- ❑ Device mobility plus varying channel quality
- ❑ Separation and merging of networks possible
- ❑ Asymmetric connections possible



time =  $t_1$



time =  $t_2$

—————> good link  
 —————> weak link



# Traditional routing algorithms

## Distance Vector

- ❑ periodic exchange of messages with all physical neighbors that contain information about who can be reached at what distance
- ❑ selection of the shortest path if several paths available

## Link State

- ❑ periodic notification of all routers about the current state of all physical links
- ❑ router get a complete picture of the network

## Example

- ❑ ARPA packet radio network (1973), DV-Routing
- ❑ every 7.5s exchange of routing tables including link quality
- ❑ updating of tables also by reception of packets
- ❑ routing problems solved with limited flooding



# Routing in ad-hoc networks

THE big topic in many research projects

- ❑ Far more than 50 different proposals exist
- ❑ The most simplest one: Flooding!

## Reasons

- ❑ Classical approaches from fixed networks fail
  - Very slow convergence, large overhead
- ❑ High dynamicity, low bandwidth, low computing power

## Metrics for routing

- ❑ Minimal
  - Number of nodes, loss rate, delay, congestion, interference ...
- ❑ Maximal
  - Stability of the logical network, battery run-time, time of connectivity ...



# Problems of traditional routing algorithms

## Dynamic of the topology

- ❑ frequent changes of connections, connection quality, participants

## Limited performance of mobile systems

- ❑ periodic updates of routing tables need energy without contributing to the transmission of user data, sleep modes difficult to realize
- ❑ limited bandwidth of the system is reduced even more due to the exchange of routing information
- ❑ links can be asymmetric, i.e., they can have a direction dependent transmission quality



# DSDV (Destination Sequenced Distance Vector)

## Early work

- ❑ on demand version: AODV

## Expansion of distance vector routing

## Sequence numbers for all routing updates

- ❑ assures in-order execution of all updates
- ❑ avoids loops and inconsistencies

## Decrease of update frequency

- ❑ store time between first and best announcement of a path
- ❑ inhibit update if it seems to be unstable (based on the stored time values)





# Dynamic source routing I

Split routing into discovering a path and maintaining a path

## Discover a path

- ❑ only if a path for sending packets to a certain destination is needed and no path is currently available

## Maintaining a path

- ❑ only while the path is in use one has to make sure that it can be used continuously

No periodic updates needed!



# Dynamic source routing II

## Path discovery

- ❑ broadcast a packet with destination address and unique ID
- ❑ if a station receives a broadcast packet
  - if the station is the receiver (i.e., has the correct destination address) then return the packet to the sender (path was collected in the packet)
  - if the packet has already been received earlier (identified via ID) then discard the packet
  - otherwise, append own address and broadcast packet
- ❑ sender receives packet with the current path (address list)

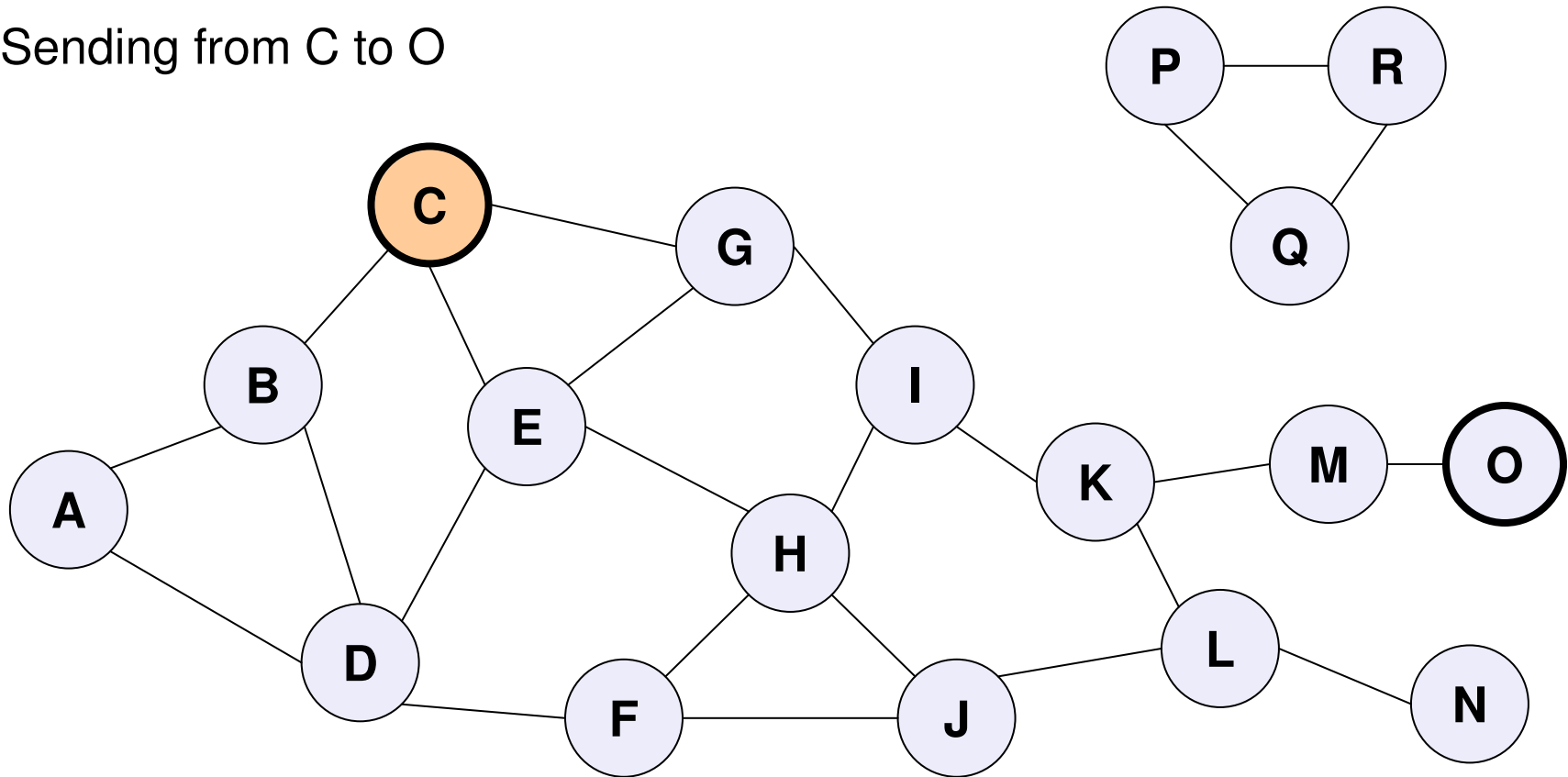
## Optimizations

- ❑ limit broadcasting if maximum diameter of the network is known
- ❑ caching of address lists (i.e. paths) with help of passing packets
  - stations can use the cached information for path discovery (own paths or paths for other hosts)



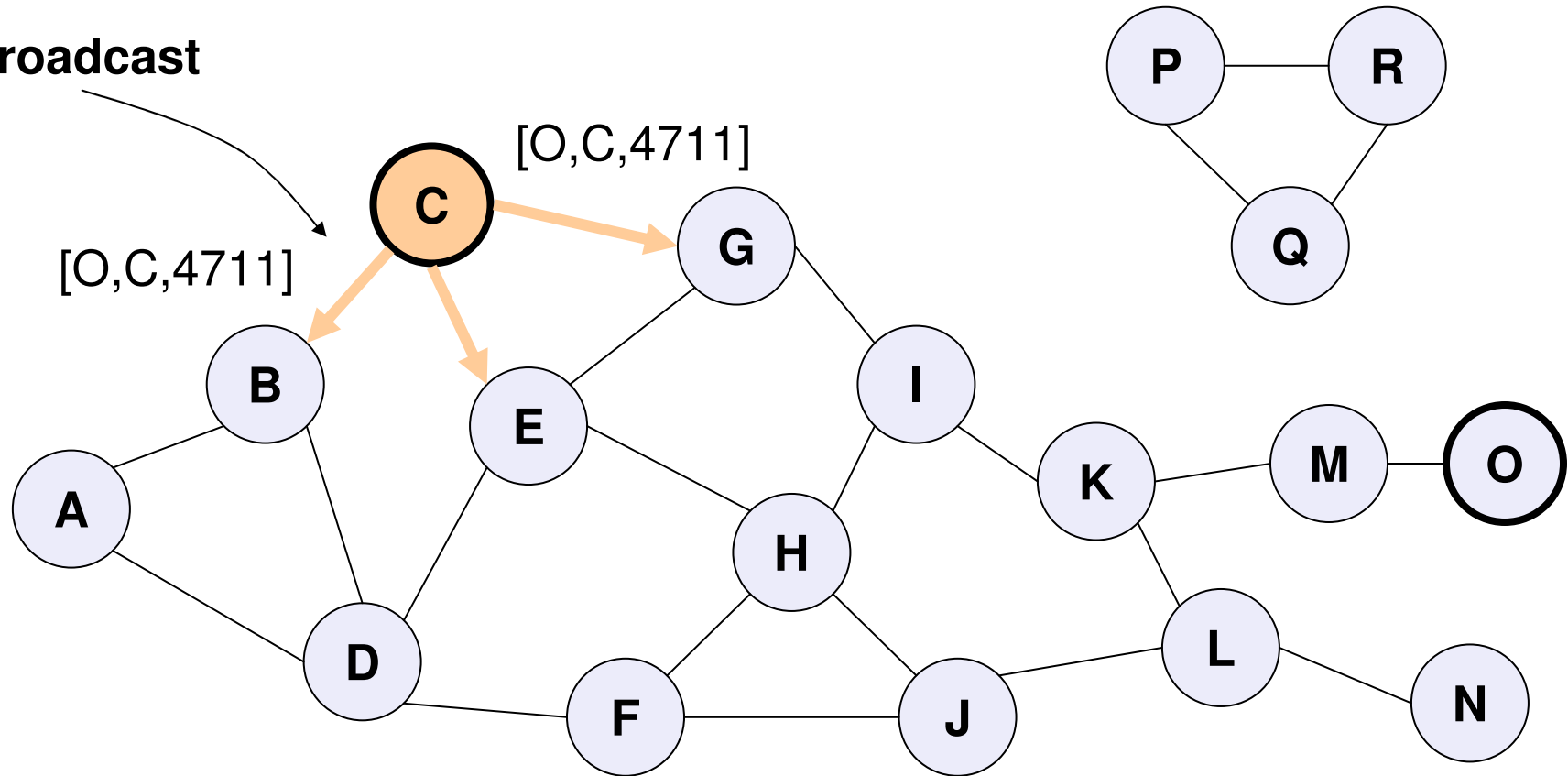
# DSR: Route Discovery

Sending from C to O

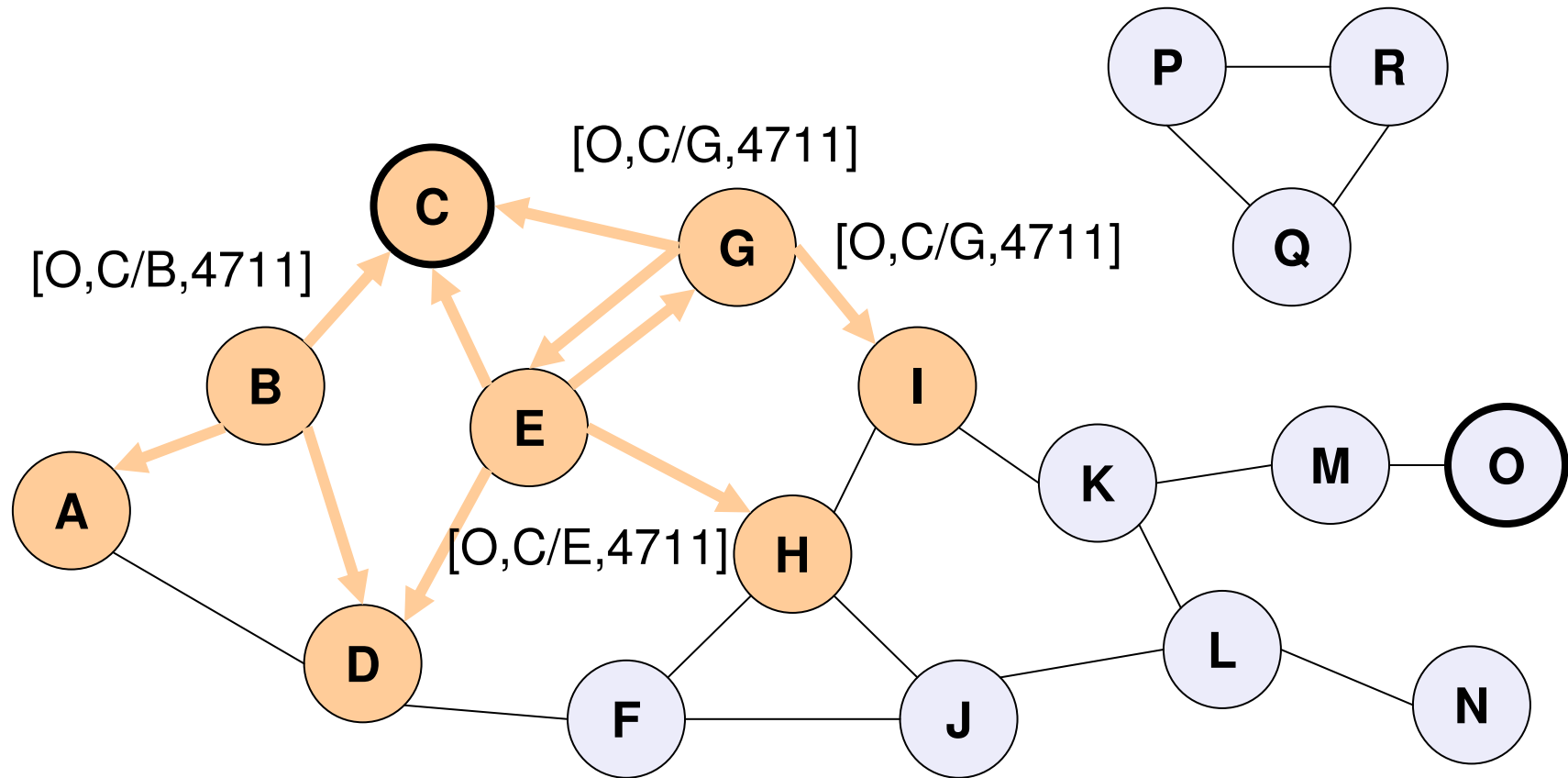


# DSR: Route Discovery

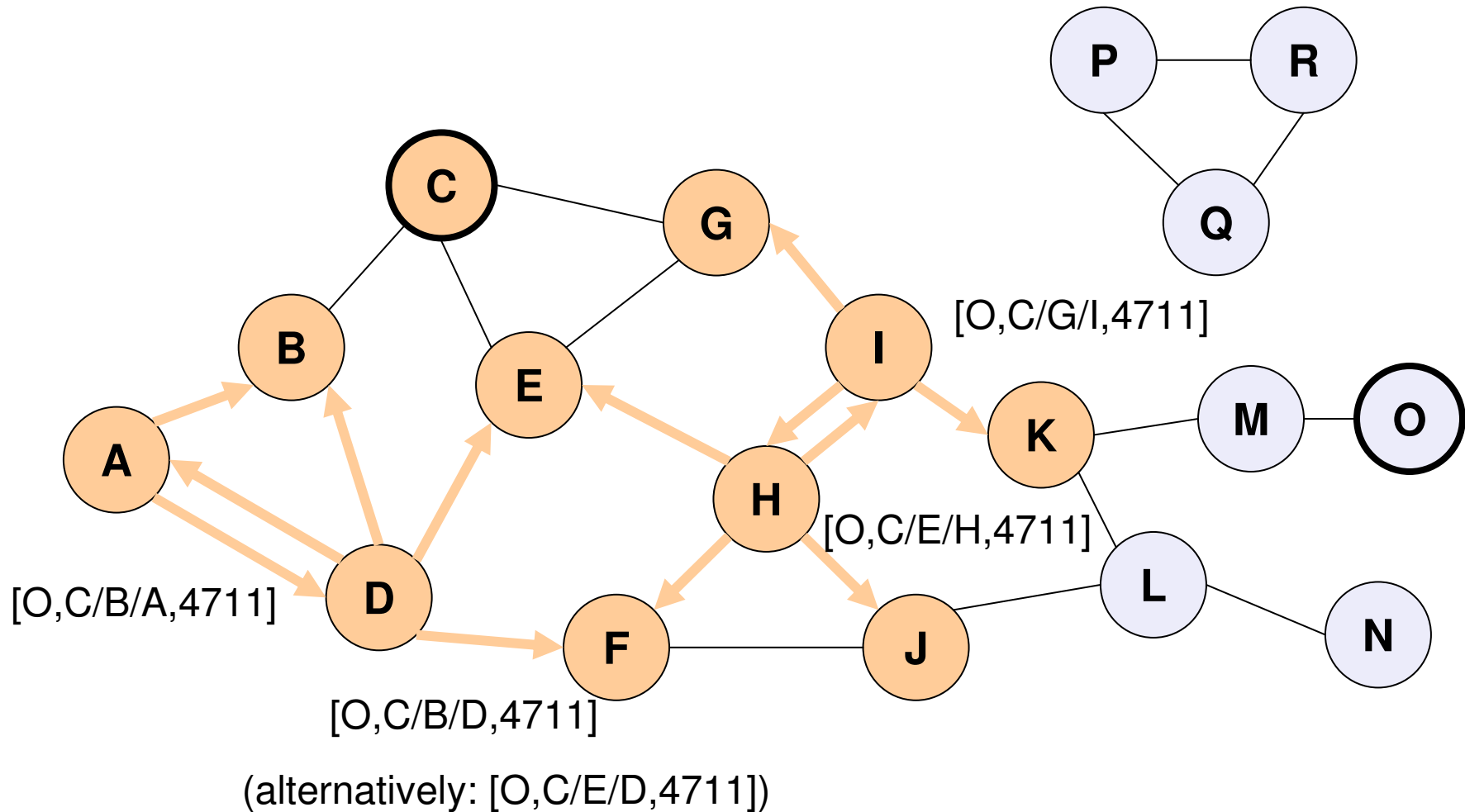
**Broadcast**



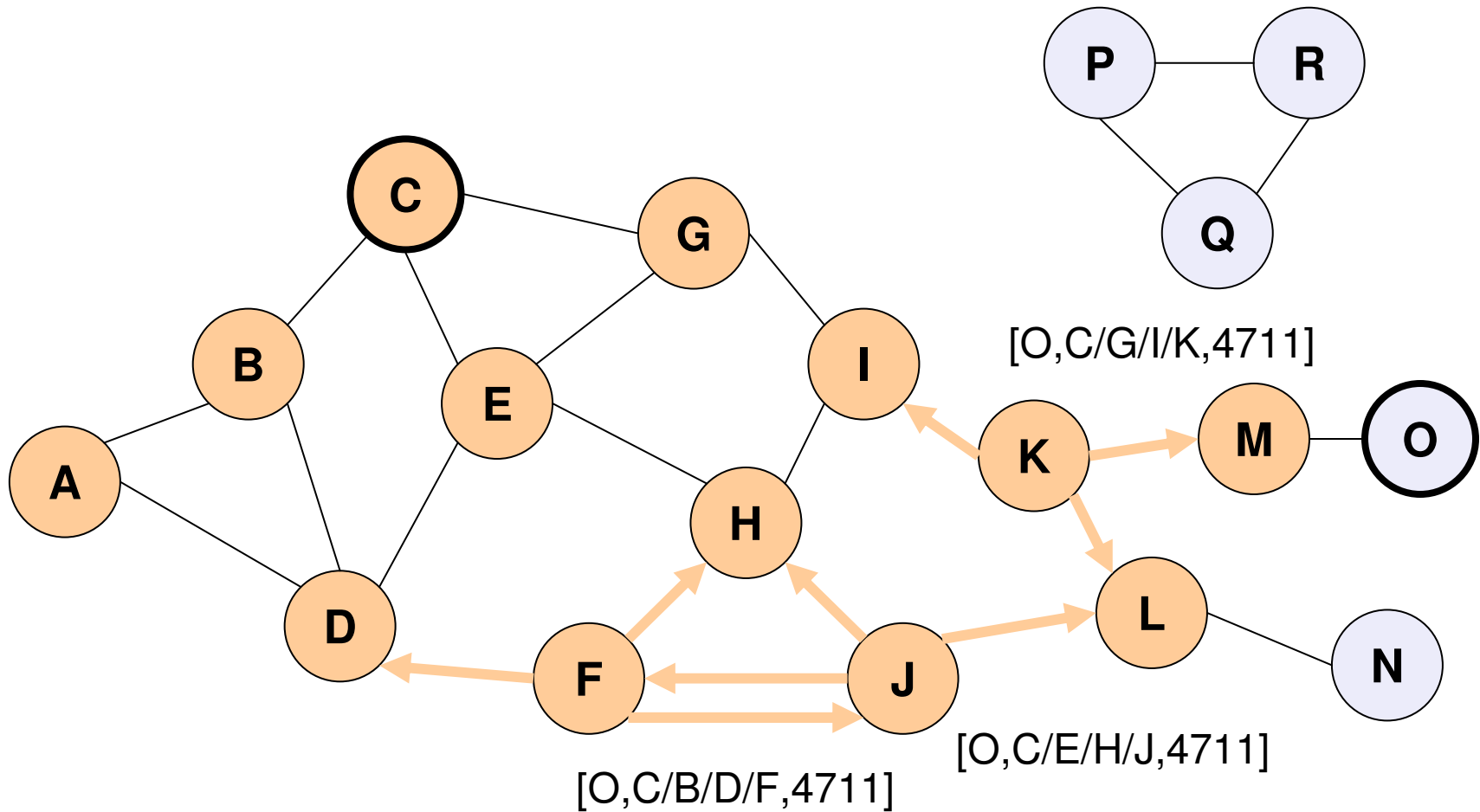
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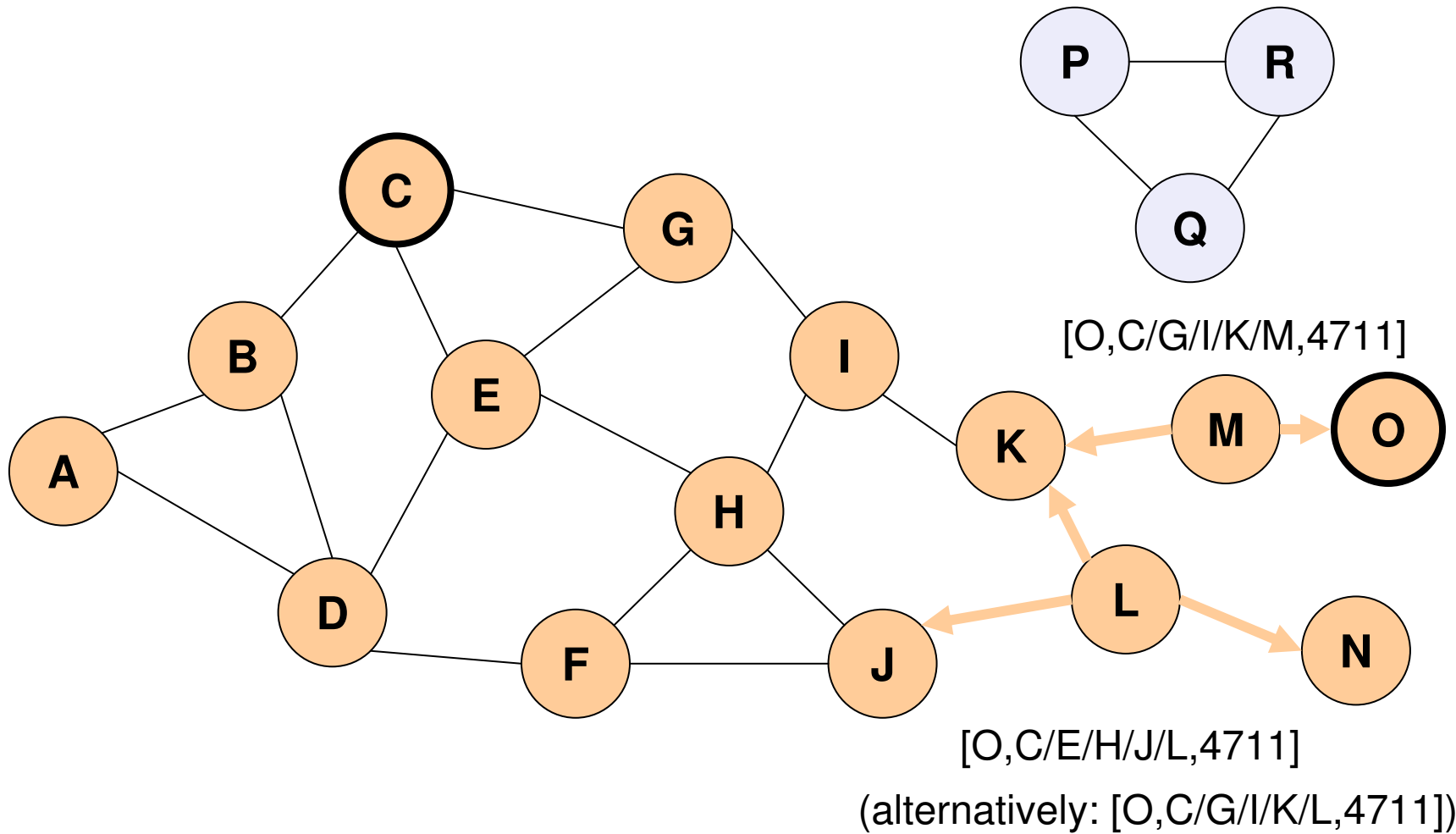
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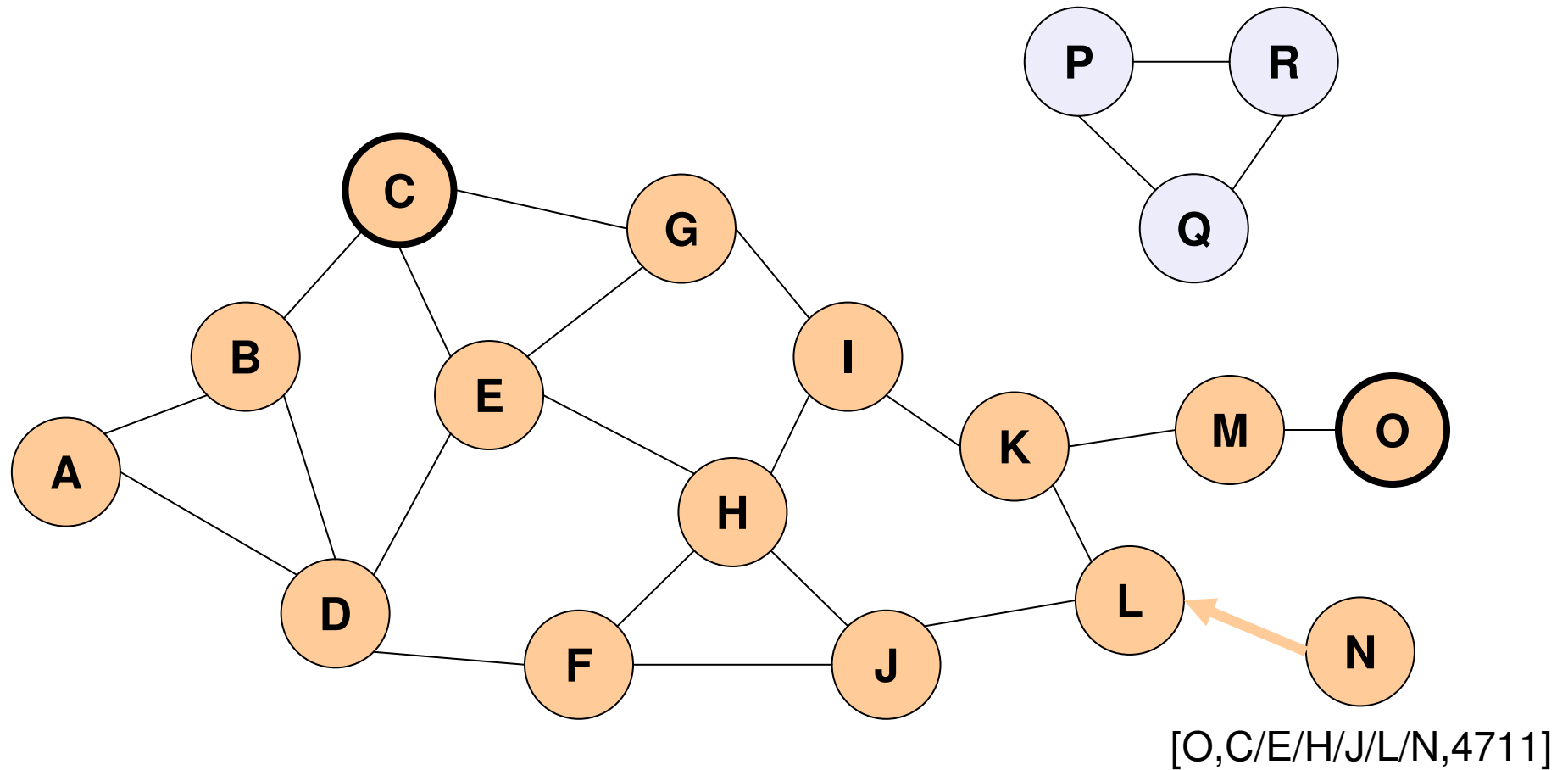


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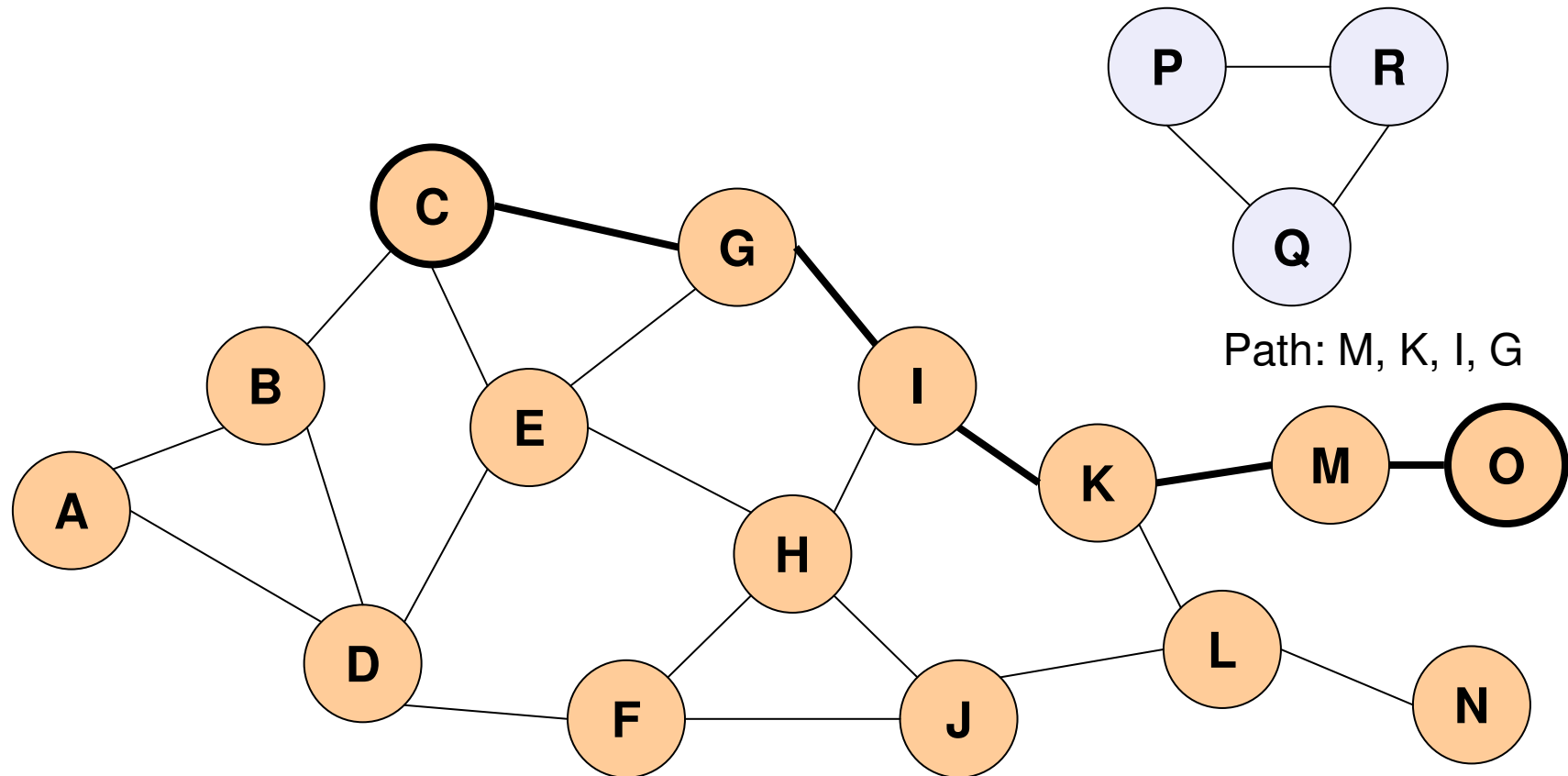




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# Dynamic Source Routing III

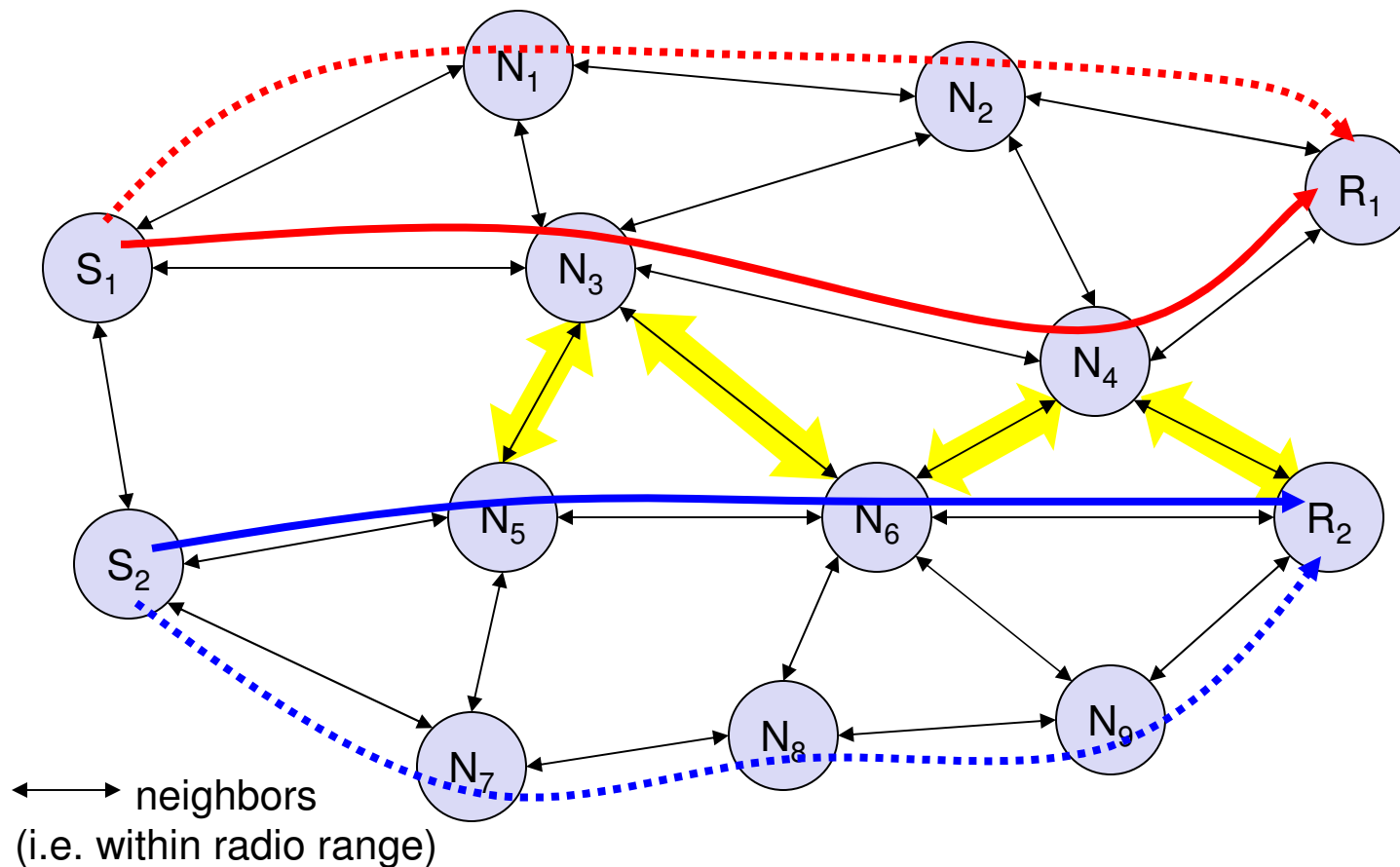
## Maintaining paths

- ❑ after sending a packet
  - wait for a layer 2 acknowledgement (if applicable)
  - listen into the medium to detect if other stations forward the packet (if possible)
  - request an explicit acknowledgement
- ❑ if a station encounters problems it can inform the sender of a packet or look-up a new path locally



# Interference-based routing

Routing based on assumptions about **interference** between signals



# Examples for interference based routing

## Least Interference Routing (LIR)

- ❑ calculate the cost of a path based on the number of stations that can receive a transmission

## Max-Min Residual Capacity Routing (MMRCR)

- ❑ calculate the cost of a path based on a probability function of successful transmissions and interference

## Least Resistance Routing (LRR)

- ❑ calculate the cost of a path based on interference, jamming and other transmissions

LIR is very simple to implement, only information from direct neighbors is necessary



# A plethora of ad hoc routing protocols

## Flat

- ❑ proactive
  - FSLS – Fuzzy Sighted Link State
  - FSR – Fisheye State Routing
  - OLSR – Optimised Link State Routing Protocol
  - TBRPF – Topology Broadcast Based on Reverse Path Forwarding
- ❑ reactive
  - **AODV** – Ad hoc On demand Distance Vector
  - DSR – Dynamic Source Routing

## Hierarchical

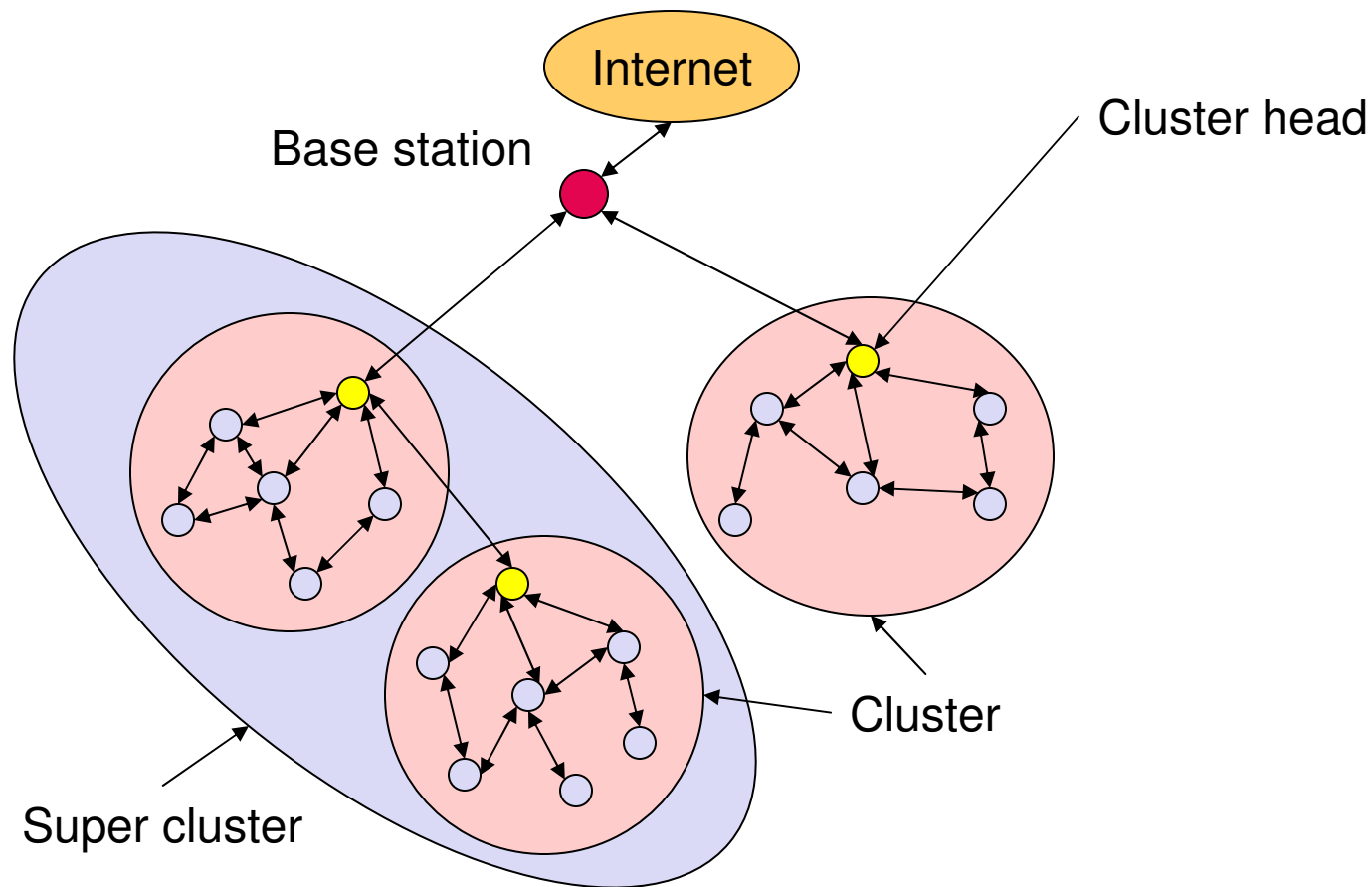
- ❑ CGSR – Clusterhead-Gateway Switch Routing
- ❑ HSR – Hierarchical State Routing
- ❑ LANMAR – Landmark Ad Hoc Routing
- ❑ ZRP – Zone Routing Protocol

## Geographic position assisted

- ❑ DREAM – Distance Routing Effect Algorithm for Mobility
- ❑ GeoCast – Geographic Addressing and Routing
- ❑ GPSR – Greedy Perimeter Stateless Routing
- ❑ LAR – Location-Aided Routing



# Clustering of ad-hoc networks



# Further difficulties and research areas

## Auto-Configuration

- ❑ Assignment of addresses, function, profile, program, ...

## Service discovery

- ❑ Discovery of services and service providers

## Multicast

- ❑ Transmission to a selected group of receivers

## Quality-of-Service

- ❑ Maintenance of a certain transmission quality

## Power control

- ❑ Minimizing interference, energy conservation mechanisms

## Security

- ❑ Data integrity, protection from attacks (e.g. Denial of Service)

## Scalability

- ❑ 10 nodes? 100 nodes? 1000 nodes? 10000 nodes?

## Integration with fixed networks

